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<p>(54) Title: MASS-STORAGE APPLICATIONS OF LOCAL PROBE ARRAYS</p> <p>(57) Abstract</p> <p>The present invention concerns a storage device comprising a local probe array (72) and a storage medium (70) with an array of storage fields (71). The local probe array (72) is situated opposite to said storage medium (70) such that each local probe (73) of the local probe array can be scanned over the corresponding storage field (31, 40, 71). The storage device further comprises means (RF, 44; CF, 45) for distinguishing between information (41) to be erased from a first section of the storage medium and information in this section which is not to be erased, means for selectively copying said information which is not to be erased into a memory, preferably another section (32) of said medium, and means for erasing the whole first section.</p> <p>40 41 42 43 44 45 RF CF 70 71 72 73 74 75.1 75.2 76.x level 3 level 2 level 1</p>			

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DESCRIPTION**Mass-Storage Applications of Local Probe Arrays**

5

TECHNICAL FIELD

The present invention relates to storage devices based on local probe
10 arrays and appropriate storage media.

BACKGROUND OF THE INVENTION

Improvements in semiconductor processing techniques have led to drastic
15 reductions in the size of todays computers. However, while the
microprocessors, displays and other components are getting cheaper and
smaller, the bulk data memory units limit the overall size reduction. For
further reduction in size and power consumption the conventional disk drive
storage systems need to be replaced by small, high capacity storage
20 devices. Such disk drive storages which usually have a rotating memory
store as much as one Gigabyte of data, but offer modest read/write speed
on the order of milliseconds only. Solid state memories, on the other hand,
offer much higher read/write speeds on the order of nanoseconds. However,
their storage capacity is limited to a few Megabit. In terms of cost per bit,
25 rotating memory is cheaper, but it is also much more mechanically
unreliable because of the moving parts.

There is a clear demand for single storage devices having storage capacity
of more than 1 Terabit. It is further important for such a storage device, in
30 particular when being used in a multimedia system where image frames
need to be retrieved in a fast and consecutive manner, that very high data
rates (read/write speed) can be achieved. Other important aspects are
power consumption, overall weight and size, reliability, data security, and

1 shock resistance (if used in portable computer systems). With a storage device which combines the capacity of a rotating memory with the speed, size, power consumption and reliability of solid state memories, computers would take another quantum leap in performance and compactness.

5

The development of scanning tunneling and atomic force microscopes has led to first storage systems which make use of parallel local probes.

10 A scanning tunneling storage system has been proposed in the European patent EP 247219. This system comprises current detectors being attached to an array of cantilevers. A storage medium is placed opposite to the array. The storage medium is displaced by means of a two-dimensional piezoelectric positioning device. There is no adequate approach for erasing of information disclosed.

15

20 In US patent 5,307,311 a memory device is described which makes use of a very large set of independently operating subdevices. It employs an array of hundreds of microcantilevers having an area in which bits are stored. Opposite to these cantilevers there are hundreds of read/write heads which are similar in nature to scanning tunneling or atomic force microscope scanning tips. Each cantilever is moved in an oscillatory manner such that the respective read/write head scans over the bits stored thereon.

25 It is a clear disadvantage of such scanning probe storage systems that they are very complex. Each such substorage needs its own driving mechanism and electric wiring as well as write and read lines connected to the numerous read/write heads.

30 It is essential for a storage device that information can be erased. In particular in case of scanning probe storage systems which have been proposed so far no reliable and satisfying erase technique has been proposed. Recent material investigations have revealed special materials which are in principle suitable as erasable storage medium. However, steps

1 required to erase information being stored in such a material are either too slow, or cannot be controlled properly to facilitate erasure of single bits within a storage medium.

5 In view of the disadvantages of known media suitable for use in scanning probe storage systems there is a need for improved materials and storage concepts in particular to overcome the problems inherent to the erase process. The high resolution available through the application of a STM or AFM is certainly a most desirable attribute. However, for purposes of mass

10 data storage, the erase capabilities of known approaches are much too slow for any practical application.

It is an object of the present invention to provide a method and apparatus which improves known scanning probe storage systems.

15 It is a further object of the present invention to provide a method and apparatus which allows fast and reliable erasure of stored information.

SUMMARY OF THE INVENTION

20 This has been achieved by the provision of a storage device having

- a storage medium (10) on which information can be stored in form of perturbations (12),
- a local probe array (20) facing said medium (10) such that its probes (14) scan said medium (10),
- means for obtaining signals from each of said probes (READ) when scanning over said perturbations, and
- means for writing perturbations onto said medium (WRITE).

1 Said storage device is characterized in that it further comprises

5 • means for distinguishing between information to be erased from a first section of said medium and information in this section which is not to be erased,

10 • means for selectively copying said information which is not to be erased into a memory, preferably another section of said medium, and

15 • means for erasing the whole first section (ERASE).

DESCRIPTION OF THE DRAWINGS

15 The invention is described in detail below with reference to the following schematic drawings:

20 **FIG. 1A** shows a perspective view of part of a storage device, in accordance with the present invention.

25 **FIG. 1B** shows a cross-section of the storage device of Figure 1A.

30 **FIG. 2** shows a perspective view of the storage device, part of which are illustrated in Figures 1A and 1B, in accordance with the present invention.

35 **FIG. 3** is a top view of a storage medium structured in accordance with the present invention.

40 **FIG. 4** is a top view of a storage field, in accordance with the present invention.

1 **FIG. 5A** shows a cross-section of part of a storage device, in
accordance with the present invention.

5 **FIG. 5B** is a top view of the storage device of Figure 5A, in accordance
with the present invention.

10 **FIG. 6** shows a top view of a storage field and a table maintained in
another field, in accordance with the present invention.

15 **FIG. 7** shows a perspective view of another storage device comprising
of at least 3 levels, in accordance with the present invention.

20 **FIG. 8** is a top view of a part of a storage medium comprising
sub-arrays, in accordance with the present invention.

25 **FIG. 9** is a schematic block diagram used to explain the multiplexed
addressing scheme, according to the present invention.

30 **FIG. 10** is a schematic block diagram of the read-out electronic,
according to the present invention.

1

GENERAL DESCRIPTION

5 Before different embodiments of the present invention are described, the basic elements of the storage devices in accordance with the present invention are addressed.

Cantilevers and cantilever arrays:

10 Cantilevers are well known elements which are easy to make. Existing semiconductor fabrication processes can be employed. In essence, the techniques of micromachining are employed to create discrete cantilevers and arrays of cantilevers. When dimensioning such cantilevers, one has to take into account specific parameters of the material used as substrate in which the cantilevers are formed. When properly designing such an array, it can be mass-produced at low cost and with high yield.

15

Usually, cantilevers and cantilever arrays are made by etching away portions of a silicon substrate. This substrate is normally (100) or (111) oriented. (100) oriented silicon could for example be wet etched using ethyl diamine pyrocatechol or KOH solutions. Wet etching techniques are 20 generally dependent on crystallographic orientation of the substrate, e.g. (100) oriented silicon shows a very low etch rate of the (111) plane, leading to a good etch stop along the (111) axis which generates well defined etch planes with 54.7 ° angles from (100). An alternative approach makes use of dry etching techniques, e.g. reactive-ion beam etching (RIE), chemically 25 assisted ion beam etching, or microwave assisted plasma etching. Depending on process conditions, deep and anisotropic structures can be obtained leading to excellent dimensional control. Masks can be employed to define the structures to be etched. The cantilevers used can have any shape that can be obtained by photolithography and etching. The 30 cross-sectional shape could for example be rectangular, round, elliptical, or polygonal.

1 Also suited for the fabrication of cantilevers are other semiconducting
2 materials, like gallium arsenide, as reported in "dynamic Micromechanics on
3 Silicon: Techniques and Devices", K.E. Petersen, IEEE Transactions on
4 Electronic Devices, Vol. ED25, No. 10, 1978, pp. 1241 - 1249.

5

Local probes and local probe arrays:

Usually a tip is used as local probe. Different techniques are known to produce such tips. They can for example be made by isotropic wet or dry etching in combination with the oxidation of a single crystal material such as 10 silicon. The following materials are well suited for making local probes and local probe arrays: tungsten, tungsten alloy, platinum, molybdenum, silicon (doped or undoped), doped diamond, any refractory metal, or conductive ceramics, to name some. The combination of wet or dry etching and liftoff plus oxidation leads to very sharp pointed cones. The sharper the tips are, 15 the denser information on a storage medium can be stored, i.e. the higher the storage capacity of a storage device will be. The probes can be coated with an appropriate metal such as gold, for example. In US patent 5 204 581 it is described in detail how to make tips or arrays of tips which can be used in connection with the present invention. Another example for 20 the microfabrication of a tip is disclosed in the article "Silicon cantilevers and tips for scanning force microscopy", J. Brugger et al., Sensors and Actuators A, Vol. 34, 1992, pp. 193 - 200. It is important to note that by means of batch fabrication local probe arrays can be made in a reproducible and cheap manner.

25

Driving circuitry:

Certain means, including driving circuitry, preamplifiers, and an appropriate wiring for reading and writing information need to be provided. To make these means one can employ existing tools and processes common to the 30 semiconductor and solid-state industries. The driving electronics as well as the probes call for circuitry like that used in scanning tunneling microscopy (STM) and atomic force microscopy (AFM) systems, albeit shrunk to extremely small size. The miniaturization is mandatory to obtain short

1 interconnections, high speeds, and less power for circuitry. To cope with the high data rates which can be achieved by a highly parallel scanning probe array, one needs to provide very fast electronic circuitry for reading and writing information.

5

Storage media:

A storage medium in accordance with the present invention is partitioned into storage fields in which the information is actually stored. The media which can be used in connection with the present invention can be grouped 10 as follows. In general, perturbations can be formed in a medium, or removed therefrom by locally

- creating or altering the topographic features or composition;
- 15 • altering the crystalline phase;
- creating or destructing electronic states;
- filling or emptying existing electronic states;
- 20 • creating or altering domain structures or polarization states;
- creating or altering chemical bonds.

25 In addition to the above examples, any combination of physical or chemical effects can be used. A good and detailed description of the different media suited is given in US patent 5 307 311.

Another approach, not explicitly mentioned in the above US patent, would 30 be to use a very soft, wax-like material, polymer, or liquid crystal in which perturbations are created either by locally heating the material, or by imprinting patterns or pits by moving the probe up and down. By heating up the material so as to melt it locally, or over a larger area, e.g. one

1 storage field, one could clean it up (erase). The heating could be achieved
electrically, e.g. by local heating element such as resistors, or by means of
a laser beam. It is for example possible to provide each local probe array
with a heat source, e.g. a resistor, such that the warm probe generates
5 perturbations in the storage medium. By means of heating elements
integrated into said storage medium, or placed on the backside thereof,
whole storage sections can be erased at once.

10 Perturbations can also be generated by employing the tunneling effect in
order to move and remove atoms. This approach, according to which the
information is stored in the form of patterns of atoms adsorbed on the
surface of a storage medium, is described in US patent 4 575 822. The
probe is maintained at tunneling distance from the medium to remove
15 individual atoms from the medium for writing, and to detect variations of the
tunneling current caused by the presence or absence of atoms in scanned
locations for reading.

20 Well suited are nitride-oxide-silicon media, as shown in connection with the
first embodiment of the present invention. Either a scanning tunneling tip, or
a scanning capacitance probe placed at a certain distance with respect to
the nitride-oxide-silicon medium can be used to measure and to detect
perturbations. Further details are given in the article "Charge storage in a
nitride-oxide-silicon medium by scanning capacitance microscopy",
R.C. Barrett et al., Journal of Applied Physics, Vol. 70, No. 5, 1 Sept. 1991,
25 pp. 2725 - 2733.

30 In addition to permanent storage media, temporary storage media are
known. One example is a medium in which electric charges are trapped
when writing information into it. After a while this information disappears
because the electron charges get spread out. A material which shows such
a behavior could also be used in connection with the present invention, as
will be described later.

- 1 Beneath STM based approaches, AFM techniques which are capable of seeing contrasts (perturbations) on insulating and magnetic storage media without using any electrical currents, or any combinations of AFM and STM can be used.
- 5 The different storage media can be structured so as to provide separate storage fields thereon. These fields are arranged such that they do not overlap. Usually there is a certain distance left between neighboring fields to ensure that the data of the respective fields are clearly distinguishable.
- 10 The more densely these storage fields can be arranged, the higher the storage density of the overall storage device are.

Scan movement:

- 15 In the present invention different approaches for achieving a scan movement of the probes with respect to the storage medium are made use of. The simplest approach is to move either the whole local probe array while the storage medium's position remains unchanged, or the other way around. An additional degree of freedom can be achieved if the storage medium and the local probe array are moved. It is, for example,
- 20 advantageous to slowly move the storage medium back and forth in a direction parallel to a first axis. The scan excursion of the storage medium is chosen such that the local probes do not exceed the corresponding storage fields, i.e. the scan excursion is approximately equal to the dimension of the storage fields. At the same time the local probe array is
- 25 step-wise moved perpendicular. Due to the combined movement of the local probe array and the storage medium, a first row of the storage fields is scanned. Then, the probes jump to the next row before the storage medium moves back. This next row is now scanned in the opposite direction, and so forth. This approach is known as 'basket-weave scanning' in the area of
- 30 scanning electron microscopy.

One can use a variety of scanning schemes including pulsed and continuous scanning as well as varying speeds.

1 **Tracking:**

In a storage system where the local probes do not have their own actuators, tracking can be achieved by providing special tracks at the edges of the storage medium which are watched and followed by the moving local probe array. Tracking can for example be achieved by mechanical means which guide the local probe array along a predefined direction. Such mechanical tracking means can be combined with, or completely replaced by, optical or other contactless tracking means. Deflection sensors on the local probe array, which are arranged such that they interact with counterparts on the storage medium, allow to precisely move the array parallel to the storage medium. Special tracking marks can be used to mark the boundaries of the storage fields and/or storage medium. If such a tracking mark is reached the local probe array can be triggered to move backwards, or to jump to another position, for example.

15

Tracking is also important to ensure that the storage device will return into a well defined position after power-off or after it has been parked.

In addition to the actuators employed to move the whole local probe array and/or the storage medium, each local probe or the cantilever on which it is situated might have its own actuator. Each such actuator needs appropriate wiring and driving circuitry. Two different schemes are conceivable if one employs such semi-autonomous local probes: 1) the local probe actuators can be used to compensate fabrication tolerances and to level out local surface roughness of the storage medium; 2) The local probes themselves might contribute to the scan movement in which case more complex local actuators are required. If such local probe actuators are employed, one might also envision to provide tracking on a storage field basis. One example would be to arrange a long chain of atoms along a row such that these atoms can be detected and followed by the local probe.

1 **Actuators:**

Each of the cantilevers might comprise an actuator so as to displace it from a relaxed position to a deflected position or vice versa. The displacement achieved by an actuator might be damped to prevent damage, for example.

5

The different cantilever movements which can be achieved with an integrated piezoelectric actuator, also referred to as transducer, are depicted in Figures 34 - 37 of the published PCT patent application WO 89/07256. The variety of movements which can be achieved facilitates local 10 probe storage devices where each of the cantilevers out of an array of cantilevers has additional freedom of action. It is conceivable that the array as a whole is scanned over the storage medium in a slow, closed loop manner and that each cantilever scans a smaller subarea which is currently 15 in reach.

15

Coarse Actuators:

In addition to the actuator of each cantilever, a coarse macroscopic distance adjustment might be required, depending of the particular application. A coarse actuator might be employed to move the probes of an array within 20 reach of the storage medium and the fine actuator means of the cantilevers then provide the operating gap control. Coarse displacement can be used to compensate manufacturing tolerances, to move the array in a park position when not being used, or while a certain area is erased in accordance with the present invention. For coarse displacement PZT (a 25 piezoelectric ceramic material; Lead Zirconate Titanate) actuators or precision levers and micrometer screws can be used.

The above mentioned actuators need specific driving circuitry which can either be integrated into the present storage device, or which can be carried 30 out separately. It is important that the scan movement is synchronized with the multiplexers/demultiplexers used for reading and writing of data.

1 Interface electronics:

In addition to the driving means and read/write electronics kind of an interface circuitry is needed. Such a circuitry might include a microprocessor, multiplexer/demultiplexer, parallel-to-serial converters and serial-to-parallel converters, digital/analog and analog/digital conversion circuits and so forth. Of particular importance are error correction means. For some applications it is advisable to employ a microprocessor which coordinates all activities of the storage device. It is likewise conceivable to design the present storage device such that it cooperates with the microprocessor of a computer of which the storage device is a part. Very fast electronic or optical switches might be employed to cope with the huge amount of data accessible through hundreds or thousands of local probes of the present storage devices.

15 Before a complete storage device in accordance with the present inventions is described, the basic building blocks thereof are now defined with reference to Figure 1A.

Referring to Figures 1A, 1B and 2, a first embodiment of the present invention will now be described. A storage field 11.1, as shown in Figure 1A, comprises small perturbations in the storage medium 10. Each of these perturbations represents a bit. Such perturbations can be very small. Perturbations having a size of a few nm were reported in the art. To ensure that such perturbations can be easily detected, the centers of neighboring perturbations should be about 30nm apart. A storage field in accordance with the present invention might have a size of $30 \mu\text{m} \times 30 \mu\text{m}$ leading to a storage density per storage field of 10^8 bits. In Figure 2, a sub-array of four storage fields 11.1 - 11.4 is shown. 4×10^6 bits can be stored in such a sub-array. Depending on the storage medium and the process used for the detection of perturbations, a bit scan speed (read-out speed) of about 100kHz per tip can be achieved. In accordance with the present invention, a local probe array with four cantilevers 13.1 - 13.4 with the respective tips is scanned as a whole over the corresponding storage fields 11.1 - 11.4 and

1 the data in each storage field are addressed quasi-simultaneously. In
addition to this lateral movement a displacement of the array in the
direction perpendicular to the medium might be useful, e.g. when parking
the array. The maximum lateral scan excursion of a tip depends on the
5 dimension of a single storage field. In the present example, the maximum
scan excursion is $\geq 30 \mu\text{m}$. Using known x-y positioning means, an access
time in the range of 1 ms to 1 μs can be achieved. This compares very
favorably with present day disc drive access times of about 10 ms. In case
10 of a storage system as illustrated in Figure 2, a total data rate of
 $4 \times 10^5 \text{ bits/s}$ can be obtained (assuming $10^6 \text{ bits/storage field}$, 4 fields and
100kHz scan speed), of course, these numbers scale with the the scan speed
and the number of storage fields.

15 Further details of the first embodiment and the materials used are given
now. The local probe array 20 comprises metallized cantilevers 13.1 - 13.4
and tips. The storage medium consists of a semiconducting substrate 10
covered by a thin oxide layer 15. This oxide layer 15 is coated with
20 patterned dielectric layers 11.x. The size and shape of these dielectric
layers 11.x define the storage fields. The cantilevers 13.x are designed such
that the tips are pressed against the storage fields. When scanning the
whole array 20 over the storage medium 10, the tips are used to measure
the electrical properties of the corresponding storage fields.

25 The local probe array of a storage device in accordance with the first
embodiment can be simply addressed by four X-Y access lines, respectively,
which enable read/write of the local probe in a chosen storage field at a bit
position within this storage field given by the momentaneous position of the
local probe array. In case of a storage devices with $N \times N$ storage fields, N^2
X-Y access lines would be needed.

30

Well suited as storage medium is a combination of silicon as substrate 10
with nitride as dielectric layers 11.x. In such a material, information can be
stored in form of charges which are trapped in the nitride layer. Charges

1 are introduced, i.e. written, by applying a bias to the respective tip. This
causes charges to tunnel through the oxide layer 15 and to become trapped
in the nitride layer 11.x. These perturbations in form of charges have been
observed to be stable over a period of several days. The perturbations can
5 be detected, i.e. read, either by using the tips of the local probes, or by
obtaining a capacitance image of the storage fields using the metallized
cantilever. The perturbations can be largely removed, i.e. erased, by
applying a reverse bias between the tip and the storage medium. The
charge trapping mechanism is not completely reversible such that always
10 some residuals of the charges are left over. In particular if one tries to erase
perturbations by applying a short reverse biased pulse to the tip, the
perturbations are not removed completely. The longer such a storage
medium is used, the worse the signal-to-noise ratio becomes when reading
the information stored. If the information would be erased the conventional
15 way, i.e. on a bit-by-bit basis, the whole sub-array in Figure 2 would either
be totally blocked, or the read/write process in neighboring fields would be
delayed due to the slow erase process. According to the present invention,
this problem has been overcome by completely cleaning-up whole storage
fields from time to time. For this purpose either the tips can be reverse
20 biased and moved slowly over the storage fields, or additional means can
be provided to remove these charges. Neighboring storage fields could for
example be cleaned-up separately, if there are separate contacts to the
storage medium's backside and if neighboring storage fields are sufficiently
isolated from each other. This allows to apply a revered bias to one storage
25 field while writing information into an adjacent field.

30 To achieve this, a data management scheme and appropriate means for its
implementation are provided. Because there neither exists a reliable and
fast erase mechanism for bit-by-bit erasure, nor there is any in sight, a
different approach is proposed. According to the present invention new
information, or information to be changed, is written into specially selected
blank storage fields or areas of such storage fields.

1 An example is illustrated in Figure 3. In this Figure, a top view of a storage
medium 30 with several storage fields 31 is depicted. Before the storage
field 31 is erased, all information in this field which is still needed is
transferred to another, empty field 32. Likewise, a conventional memory
5 device might be employed for storing such information which is still needed.

The present storage device comprises means which help to distinguish
information which is not longer needed from information which is still
needed. This can be done by means of a table maintained in a reserved,
10 separate storage field 33, or by reserved bits, e.g. at the side of each
storage field. Two different approaches are outlined in connection with
Figures 4 and 5B. In both Figures, a storage field 40 is shown which
comprises 25 bit locations, for example. Assuming that two of these bits,
see reference numeral 41, are still needed, whereas all other bits are not
15 longer needed, these two bits have to be marked appropriately. In Figure 4,
there is a column of flags, referred to as row flags (RF) 44, and a row of
flags, referred to as column flags (CF) 45. The first and second bits 43 of the
CFs 45 are raised (e.g. 'one') and the second flag 42 of the RFs 44 is raised,
too. These three flags 42 and 43 clearly identify the bits 41 as still being
20 used.

In Figure 5B, the CFs 45 are rotated and placed next to the RFs 44. As in
Figure 4, the flags of the CFs 45 point to the respective rows and the flags of
the RFs 44 point to the respective columns of the storage field 40. The
25 advantage of this arrangement is that the flags can be read concurrently
with the bits in the storage field, provided that there is an array of parallel
levers 47 probes 46, as depicted in the cross-sectional view given in
Figure 5A. An alternative approach would be to mount all these parallel
levers 47 at the end of one and the same cantilever. In this case additional
30 means have to be provided which give each probe some freedom of
movement along the z-axis in order to level out surface roughness of the
storage medium. If one assumes a 30nm bit size, there would not be
sufficient space to realize the tips 46, because the distance from tip to tip

1 would also be 30nm. To overcome this problem, the data on the storage
medium can be arranged in an interleaved manner. If one assumes a
distance between neighboring tips of 300nm the 1st, 11th, 21st, 31st and so
forth columns would be read concurrently. Then, the probe array 46 is
5 moved 30nm and the 2nd, 12th, 22nd, 32nd etc. are read.

Instead of using a separate table or column and row flags, each bit word in
the storage field could comprise a header indicating whether the
information comprised in the word is still needed.

10 The means used to indicate what part of the information in a storage field is
still needed is hereinafter referred to as erase pointers and the process
used to distinguish between information which might be used later and
information not longer used is called erase data management. In case that
15 one uses a column and row flag for each bit, much real estate on the
storage medium is wasted for these erase pointers. Depending on how the
information in a storage field is organized, less erase pointers are needed.
In Figure 6 an example of five 8-bit words within a storage field is
illustrated. The whole sub-array, not shown in this Figure, e.g. comprises 20
20 storage fields ($n = 20$) each having the capacity for storing five 8-bit words,
i.e. the overall storage device has a capacity of 800 bits. In the present
example the 11th storage field 60 ($n = 11$) is to be erased because most of
the information is obsolete. Only the fifth word in this field 60 is still
needed. In a corresponding erase pointer table 62, a flag 61 is raised which
25 indicates that the fifth word in the 11th storage field 60 ($n = 11$) should not
be erased. The fifth word is consequentially rescued by copying into another
field before the whole storage field 60 is then cleaned-up.

30 It is obvious from this and the previous examples that there are many
different approaches conceivable to determine whether information can be
erased. The longer the bit words in the storage fields are, the less pointers
in the erase pointer table or the less erase flags are required.

1 The essential steps of the erase data management process are now
described:

1. First of all, the erase pointers are checked to detect whether a particular
5 information is still needed or not, i.e. either the corresponding flags,
headers, or a table maintained separately are consulted. In case of the
examples illustrated in Figures 4 and 5, a simple AND combination of
the respective row and column flags could be generated. The result of
this AND combination is then mapped onto the information in the
10 storage field to distinguish information no longer needed from the
information which is still needed.
2. The information identified as still being needed is then transferred from
the respective storage field into another, blank storage field (field 32 in
15 Figure 3, for example). Once all information in the originating storage
field has been 'rescued', the next step is initiated.
3. The whole originating storage field, or several such fields are now
20 erased at once. Depending on the material used as storage medium,
this can be done by illumination, heating, demagnetizing, or
discharging, as explained above. While one or several storage fields
are erased, it is advantageous to park the cantilever array, or to move it
25 in a manner that the whole surface of the storage field can be accessed
without interfering with the cantilever. A coarse actuator could for
example be employed to remove the cantilever array, or likewise a
shape memory alloy (SMA) can be used to displace or flip the array.
4. A storage field which has been erased can now be marked and treated
as a blank storage field.

30

While the information not pointed at by an erase pointer is transferred into
another, blank field, an error detection mechanism might be applied to

- 1 check whether information got corrupted. Any kind of error detection mechanism, like cyclic-redundancy check, can be employed.
- 5 In addition to the pointing at information which is to be deleted, one can use a similar approach to access information which needs to be changed. Using such a 'change pointer', one can retrieve a word from a storage field. This word is then forwarded to a processor where it is processed and changed. The changed word is then stored in another, blank field and the position in the storage field where it was stored beforehand is marked by means of a 10 erase pointer, because this information is not longer needed. If the result of the calculation showed that the word does not need to be changed, no erase flag needs to be raised and there is no need to write this word into another storage field.
- 15 Since overwriting of information is not possible in the conventional sense, the storage fields of a storage device might get filled up under certain circumstances. Despite of the tremendous storage capacity this calls for a compression of the data. In accordance with the present invention a process for geometrical compression of the data, which goes hand in hand with the 20 above described erase process, is described. In the following, two different compression schemes are addressed.

In-out clean up:

As described in connection with the above erase process, a storage field is 25 searched for information not longer needed using the respective erase pointers. The information which is still needed is then transferred into another, blank storage field. According to the present invention this transfer of information is done geometrically, i.e. blank areas in originating storage field are copied and appear as blank areas in the destination storage field 30 on a bit-by-bit basis. Now, a compression mechanism (fuzzy puzzle clean-up) starts searching for files or words from other storage fields or sub-arrays that fit well into the blank gaps in the destination field. Due to the fact that there are hundreds or even thousands of different storage

1 fields, the probability to find such words or files is very high. Like in a
puzzle, a destination field is now packed until its storage capacity is almost,
or in some cases even totally used up. Kind of a 'traveling salesman'
approach, making use of a 'travelling salesman algorithm', or a genetic
5 algorithm might be used such that the storage density is optimized in an
evolutionary manner. Such a 'travelling salesman algorithm' or genetic
algorithm allow fast and efficient copying of information from one storage
field directly into another. If such a smart algorithm is not implemented, it
is advantageous to copy the information still needed into a temporary
10 storage field or external storage device. The information can then be
compressed and afterwards stored in a blank storage field, for example.

15 The data compression can be further improved or modified, for example, by
applying certain rules for the selection of a storage field which is to be
emptied. This field could be selected by comparing the amount of obsolete
data of a field with the amount in other fields. The storage field with most
obsolete data, i.e. the field where only a few bits are still needed, can be
treated first. Then the next storage field follows, and so on. The rules for
cleaning-up certain storage fields can either be user-defined, or can be
20 implemented when designing a storage device.

25 Another embodiment of the present invention is now described. To obtain a
storage which is suitable as replacement of conventional storage means, the
storage capacity needs to be sufficiently high. Such a high storage capacity
can be realized in that a storage medium 70 with many storage fields 71 is
used. In the embodiment schematically illustrated in Figure 7, the storage
medium 70 carries 1000 x 1000 storage fields 71, each storage field having
a capacity of 10^6 bits. This leads to a storage device which is capable of
storing up to 1 Terabit. To access this information, a two-dimensional local
30 probe array 72 having 1000 x 1000 cantilevers 73 and probes is provided.
This two-dimensional array 72 of cantilevers is placed at a close distance to
said storage medium 70 and a probe is assigned to each storage field
thereof. In the present example the storage device makes use of the

1 tunneling effect for reading and writing of information. In order to
individually address each one of the probes, an appropriate wiring is
needed. If the storage medium 70 is kept on ground potential by using a
common contact to its backside, there is one access line per probe required.
5 This assumption only holds if the cantilevers do not require active
positioning means, such as fine actuators for z-displacement. The real
estate on the substrate of the 2-dimensional local probe array 72 would not
be sufficient to carry the 1000^2 access lines. According to the present
invention this problem is circumnavigated in that said storage device
10 comprises another layer 74, referred to a level 3, which is flipped onto the
substrate with cantilever arrays 72. This level 3 carries the read/write
lines 75.x, and select lines 76.x and, if space permits, additional electronic
read/write circuitry, for example. The employment of read/write lines 75.x
and select lines 76.x requires a multiplexed addressing scheme which will
15 be described in detail in connection with Figure 9. The three levels of the
present storage device are only illustrated schematically in Figure 7.

If the electrical wiring cannot be realized on one side of the level 3
board 74, printed circuit boards with several layers of metallization and
20 isolation sheets can be used. By means of appropriate vias the access lines
penetrate through the board and end in a metallization pad. These pads are
carried out and arranged such that an electrical contact to the probes of the
local probe array 72 is ensured if the level 3 board 74 is flipped onto it.
Solder bonding, which is normally used to connect several printed boards,
25 is well suited to be used in connection with the present invention. The
alignment accuracy which can be achieved is sufficient to realize high
density storage devices in accordance with the present invention. In
addition to using the solder bonds for alignment and mechanical
interconnection, these bonds can serve as electrical connections between
30 the local probe array and the board 74. One might also envision capacitive
electrical coupling from level 3 to level 2.

1 The embodiment illustrated in Figure 7 comprises the following materials
and the read/write process is based on the below described approach. The
storage medium 70 comprises a ferroelectric material, e.g. PZT, which is
locally polarized by means of a probe of the local probe array 72. Well
5 suited are thin film ferroelectric materials. The storage medium retains a
local remanent polarization after the tip is moved forward to a next bit
position. A reverse polarization voltage applied would polarize the medium
in the opposite direction. This effect could be used for bistable storage
devices having a higher storage density. The local polarizations are
10 non-volatile, i.e. remain unchanged if the local probe is moved away. The
perturbations in form of a local polarizations can be detected, i.e. read, by
measuring the interaction of a field applied by a local probe and the
polarization state. This interaction leads to changes in the tunneling voltage
or current. A preferred embodiment would be to scan the tips over the
15 storage fields without contacting it. By applying the required reading and
writing voltages between the tip and the corresponding storage field
information can be stored and retrieved. In order to erase information, the
tips of the local probe array 72 are operated in the field emission mode
instead of the tunneling current mode used for reading and writing. In the
20 field emission mode larger voltages or currents can be introduced to the
storage fields. This approach is not well suited for bit-by-bit erasure of
information because it has a reduced resolution and is slower. According to
the present invention, the field emission mode is used for cleaning-up whole
fields. This process of cleaning-up is controlled and carried out by special
25 data management means, already mentioned in connection with the first
embodiment.

In Figure 8 part of another storage medium 80 is shown. This storage
medium is partitioned into several sub-arrays 81. Each such sub-array has
30 a great number of storage fields 82 (see magnified sketch) as for example
described in connection with Figures 1, 2, or 7. By means of partitioning
the storage medium in such a manner, storage devices with even higher
storage capacity can be realized. One local probe array is now assigned to

1 each of the sub-arrays 81. These local probe arrays are mechanically
independent so that some sub-arrays can be in a read/write mode while
others are parked or flipped away so that the storage fields underneath
can be cleaned-up. Each of the sub-arrays might have its own read/write
5 electronics and multiplexers/demultiplexers. Higher level data management
means are provided to control data flow between sub-array, to redistribute
incoming data if some sub-arrays are currently not available or accessible,
and so on.

10 Error detection and correction (EDC) is performed in such a storage device
by each sub-array making it capable of detecting single or double bit errors
and correcting single bit errors. This is preferably done by appending one
byte of error correcting code (ECC) to each 4 bytes of memory.

15 As already mentioned above, a storage device with $N \times N$ storage fields
and $N \times N$ local probes requires N^2 address lines. A storage device with
1000 x 1000 local probes would require 10^6 address lines would not fit onto
the surface area of the local probe array. Another scheme, which is
described by reference to Figure 9, can be used in such a situation. In this
20 Figure the tips 90 of a 2-dimensional 1000 x 1000 local probe array are
schematically depicted. Read/write lines 91.x with read amplifiers 93.x and
write amplifiers 92.x are indicated on the left hand side of the local probe
array. Additionally, this storage device comprises 1000 select lines 95.x.
According to this scheme, only N read/write lines and N select lines, i.e.
25 2 x N in total, are needed and the read/write process is carried out in a
staggered manner, i.e. short pulses are fed through the select lines 95.x.
Due to this, one row of tips is addressed after the other. In the present
example, it takes 100ns until an adjacent row of tips is addressed, as
indicated in the uppermost portion of Figure 9. It takes 100 μ s to address
30 all 1000 select lines 95.x. If one uses cap-sensing for reading information
stored on a storage medium, the tip capacity 102 could be sampled at a
rate of 100ns using a simple transistor switch 101 and a read-out amplifier
100. The respective block-diagram is given in Figure 10. A similar approach

1 is already used in current dynamic random access memories (DRAMs) with
the only difference that the DRAM capacity is replaced by the capacity of the
local probe, i.e. the tip capacity.

5 In the following a computer system comprising a storage device according
to the present invention is outlined. The storage device comprises a
plurality of sub-arrays each having its own and mechanically independent
local probe array. The sub-arrays are arranged in banks and pages for
receiving data bits to be stored and for retaining data bits from said
10 devices. In addition, the storage device comprises a plurality of read/write
lines connecting to said local probes for transferring data bits thereto and
therefrom. A connector is coupled to said read/write lines in order to
couple the storage device to the computer system. The whole storage
devices is placed inside a protective enclosure. A memory controller is in
15 addition employed to control access to said storage device and its
sub-arrays, to synchronize data exchange, and to perform a memory check
at power-on. This memory controller might either comprise the data
management means according to the present invention or might interact
with such data management means via a serial or parallel interconnection
20 bus. The present storage device might further be connected to a Static
Random Access Memory (SRAM) buffer for storing a row of data from said
storage device. The above memory controller would in this case also
interact with the SRAM to load and select it. Such a SRAM can also be used
as temporary memory for clean-up.

25

The present storage devices are very well suited for use as video storage.
For example, 64-bit words can be stored and retrieved very fast. If the
storage device is linked to the image processor via a high speed bus
systems, e.g. an optical bus, images can be processed and displayed in a
30 fast and consecutive manner.

Storage devices as described above, have an enormous storage capacity.
Depending on the computer or system in which such a storage device is

1 used, it might be sufficient, if a clean-up and data compression are carried
out once a day, during night or weekend, or if no program currently takes
access to the storage. In a device as illustrated in Figure 7, data
management means are either integrated into one of the levels 1 - 3, or are
5 carried out as separate unit. If this data management means determines
that the storage device is almost full, or that there are many fields which
comprise obsolete data, a clean-up process is initiated. Depending on the
implementation the clean-up process is started only if the data management
system realizes that neither the keyboard nor any application program is
10 active. In a portable computer the clean-up process might be started if the
computer is in a stand-by mode.

As already described, the storage medium might also be heated,
illuminated, de-magnetized or the like to erase information. To do so, either
15 the cantilever array has to be flipped away, or parked at a safe distance
such that the storage medium can be accessed from above. It is also
possible to access the storage medium from underneath to erase
information. A resistive layer could for example be attached to the storage
medium's backside. By applying a voltage to this resistive layer, the
20 temperature of the medium could be increased. Depending on the storage
medium used, such a temperature treatment leads to a removal of
perturbations.

There are some storage media known in the art where perturbations can be
25 removed bit-by-bit. It was reported that such perturbations can not be
removed totally. This leads to an increasing signal-to-noise ratio the longer
a storage medium is used. A storage-field or even global erase process can
then be employed from time to time to remove residuals on the medium
which might falsely be interpreted as information. Such a storage-field or
30 global erase process usually requires quite some energy. If the present
storage devices are used in portable computers, it is recommended to wait
until the computer is connected to an AC power supply. An erase process
can then be carried out without wasting energy from the battery.

- 1 Since the local probe array is the most expensive part of a storage device according to the present invention, an increased storage capacity can be achieved as follows, without unnecessarily increasing the cost of such a device. Each local probe of a local probe array might be assigned to
- 5 several storage sections. Assuming that there are ten storage sections assigned to each local probe of a local probe array having ten such local probes, the first ten storage sections could be accessed by these ten probes in parallel. If one of these ten sections is filled up, the whole local probe array is moved so as to access the next ten storage sections, and so on.
- 10 The read access to information stored in storage sections which are currently not acted upon by said local probe arrays is slowed down, because the whole local probe array needs to be moved over a 'larger' distance before the respective storage sections can be reached by the local probes.

15

- The present invention is well suited for use in a multiprocessor environment. In such an environment the present storage device could be partitioned into multiple autonomous sub-arrays. Each such sub-array might then be assigned to a particular processor of the multiprocessor systems.
- 20 Other sub-arrays of the storage device might be shared such that results of certain computations are made available to all processors. In these shared sub-arrays the application program can be stored also.

- 25 By means of the parallel operation of miniaturized local probes arrays, according to the present invention, storage densities in the range of Terabits and data rates of 100 Gigabit/sec can be achieved.

1 **CLAIMS**

1. **Storage device comprising**
 - 5 • a storage medium (10) on which information can be stored in form of perturbations (12),
 - a local probe array (20) facing said medium (10) such that its probes (14) scan said medium (10),
 - means for obtaining signals from each of said probes (READ) when scanning over said perturbations, and
 - 10 • means for writing perturbations onto said medium (WRITE), said device being characterized in that it further comprises
 - means for distinguishing between information to be erased from a first section of said medium and information in this section which is not to be erased,
 - 15 • means for selectively copying said information which is not to be erased into a memory, preferably another section of said medium, and
 - means for erasing the whole first section (ERASE).
- 20 2. **The storage device of claim 1, wherein said means for erasing**
 - are carried by said local probe array, and
 - are arranged such that they are capable of erasing whole sections of said storage medium at once.
- 25 3. **The storage device of claim 1, wherein said means for erasing**
 - are situated at one side of said medium whereas said local probe array is situated on the other side thereof, and
 - are arranged such that they are capable of erasing whole sections of said storage medium at once.
- 30 4. **The storage device of claim 1, wherein said local probe array (20) comprises a 1- or 2-dimensional array of cantilevers each of which**

- 1 carries one or several local probes which are either kept at a certain
 distance from said medium, or are operated in contact mode.
- 5 5. The storage device of claim 1, wherein said local probe array or said
 medium is moved in a closed loop manner such that each probe scans
 a certain section of said storage medium.
- 10 6. The storage device of claim 1 or 5, wherein said medium is partitioned
 into multiple storage fields each of which being scanned by at least one
 corresponding local probe.
- 15 7. The storage device of claim 1, comprising data management means
 which distinguish between information to be erased from a first section
 of said medium and information in this section which is not to be erased
 by analyzing pointers or flags which identify said information to be
 erased.
- 20 8. The storage device of claim 7, wherein said data management means
 compares the amount of information to be erased from a first storage
 field with the amount of another storage field to determine, based on
 predefined rules, which storage field is to be erased first.
- 25 9. The storage device of claim 7, wherein said data management means
 employs a travelling salesman or genetic algorithm.
- 30 10. The storage device of claim 7, wherein either each storage field
 comprise said flags which point at bits of said storage field which are to
 be erased, or said storage medium comprises said table in which erase
 pointers are maintained which point at bits of storage fields which are to
 be erased.

- 1 11. The storage device of claim 4, wherein said cantilevers carry several probes such that whole bit-words can be written onto said medium or read from said medium.
- 5 12. The storage device of claim 1, comprising actuators which cause a lateral scan movement of said medium with respect to said local probe array, or of said local probe array with respect to said medium.
- 10 13. The storage device of claim 1, comprising actuators which partially remove or flip away said local probe array from a section of said medium which is to be erased.
- 15 14. The storage device of claim 1, comprising a board carrying address lines for read/write access to said probes, said board being flipped onto said local probe array such that each of said address lines is electrically coupled to the corresponding probe.
- 20 15. The storage device of claim 14, wherein said board comprise read/write circuitry.
- 25 16. The storage device of any of the preceding claims, wherein said medium is divided into multiple autonomous sub-arrays, each of which has multiple storage fields and a local probe array which is independent of the local probe arrays of adjacent sub-arrays.
- 30 17. The storage device of any of the preceding claims, wherein said medium comprises a semiconducting substrate (10) covered by a thin oxide layer (15) which is coated by a dielectric layer (11.x) and wherein perturbations are created by locally applying a bias by means of a probe such that charges are trapped in said medium.
18. The storage device of claim 17, wherein said probes are conducting tips and wherein said medium comprises a back contact to said

- 1 semiconducting substrate (10) such that charges are locally introduced
 into the medium if a voltage is applied between a tip and said back
 contact.
- 5 19. The storage device of claim 1, wherein said medium comprises
 ferroelectric material in which perturbations are generated by locally
 polarizing it by means of a voltage applied to said probes.
- 10 20. The storage device of claim 19, wherein said perturbations are detected
 by interaction of a field applied to a probe and the polarization of said
 perturbation.
- 15 21. The storage device of claim 19, wherein said perturbations of a section
 of said medium or of a whole storage field are erased by applying a
 sufficiently large voltage or current to said probes or other means for
 erasing information in order to operate them in the field emission
 mode.
- 20 22. The storage device of claim 1, wherein said medium comprises a
 material in which perturbations can be formed by locally heating said
 material.
- 25 23. The storage device of claim 22, wherein said perturbations of a whole
 section can be removed by heating up the whole section.
- 30 24. The storage device of claim 1, wherein said storage medium comprises
 an organic material, preferably a polymer or wax, or liquid crystal.
25. The storage device of claim 1, comprising m local probes and n storage
 sections being assigned to each of said m local probes, i.e., the storage
 device comprises m times n storage sections.

- 1 26. A mass storage comprising several storage devices according to the preceding claims and means for interconnecting said mass storage to a computing system.
- 5 27. Method for erasing information stored in a storage device comprising
 - a storage medium (10) in which said information is stored in form of perturbations (12),
 - a local probe array (20) facing said medium (10) such that its probes (14) scan said medium (10),
 - 10 • means for obtaining signals from each of said probes (READ) when scanning over said perturbations, and
 - means for writing perturbations onto said medium (WRITE), characterized by the following steps:
 - distinguishing between information to be erased from a first section of said medium and information in this section which is not to be erased,
 - selectively copying said information which is not to be erased into a memory, preferably another section of said medium, and
 - 15 erasing the whole first section (ERASE).
- 20 28. The method of claim 27, wherein information to be erased from a first section of said medium and information in this section which is not to be erased is distinguished by analyzing pointers or flags identifying said information to be erased.
- 25 29. The method of claim 27, wherein the amount of information to be erased from a first storage field is compared with the amount of other storage fields in order to decide which storage field is to be erased first.
- 30 30. The method of claim 27, wherein information read from a storage field which is to be erased is compressed before it is written into said memory.

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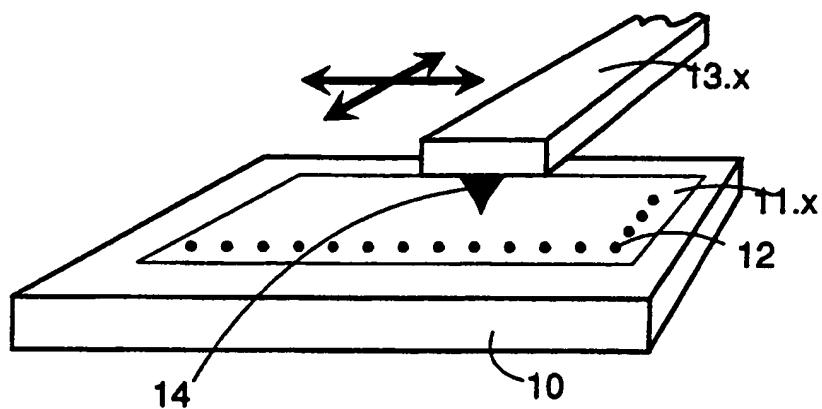


FIG. 1A

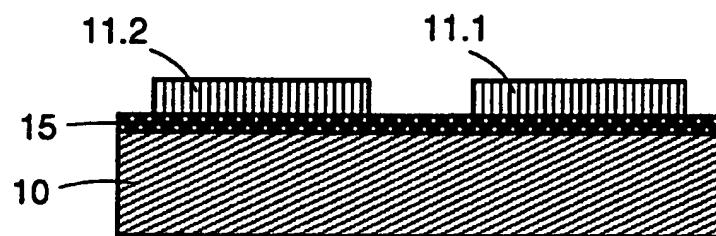


FIG. 1B

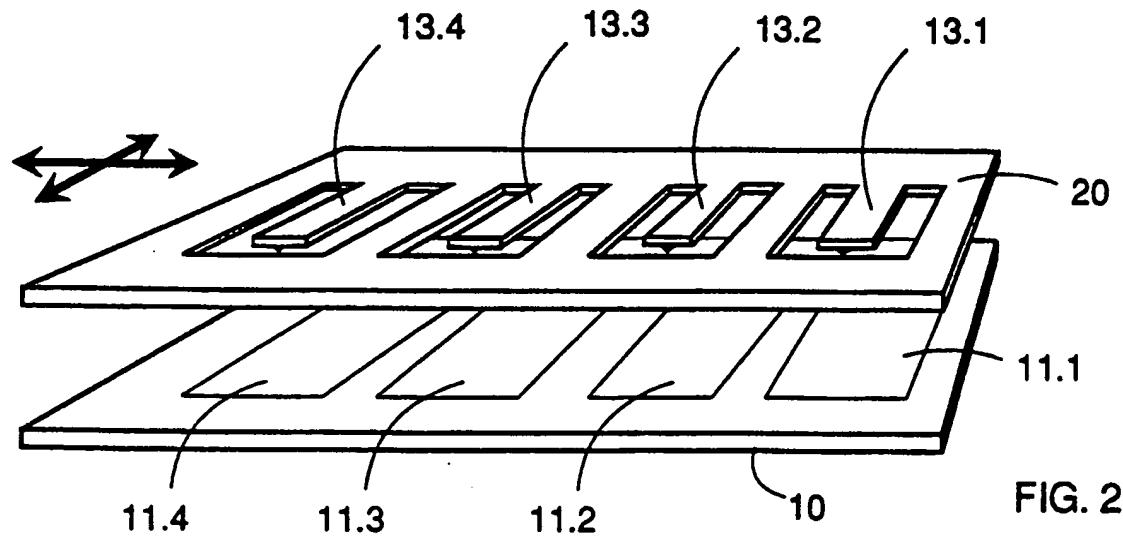
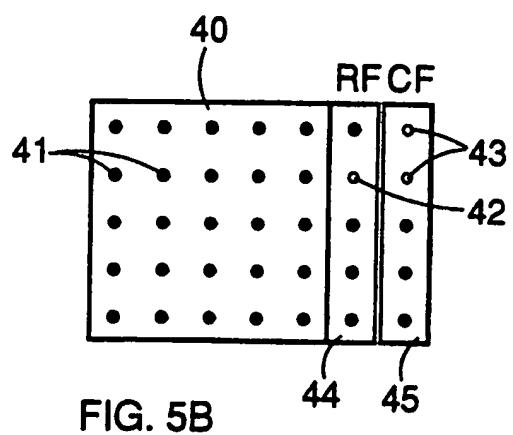
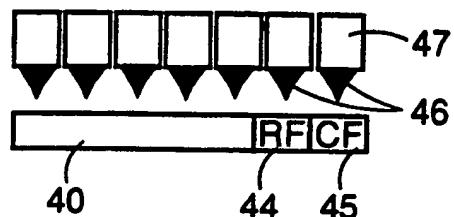
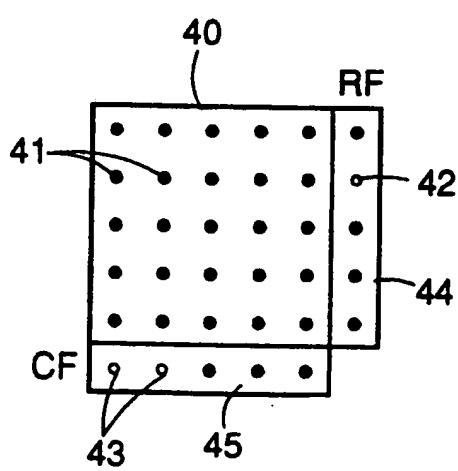
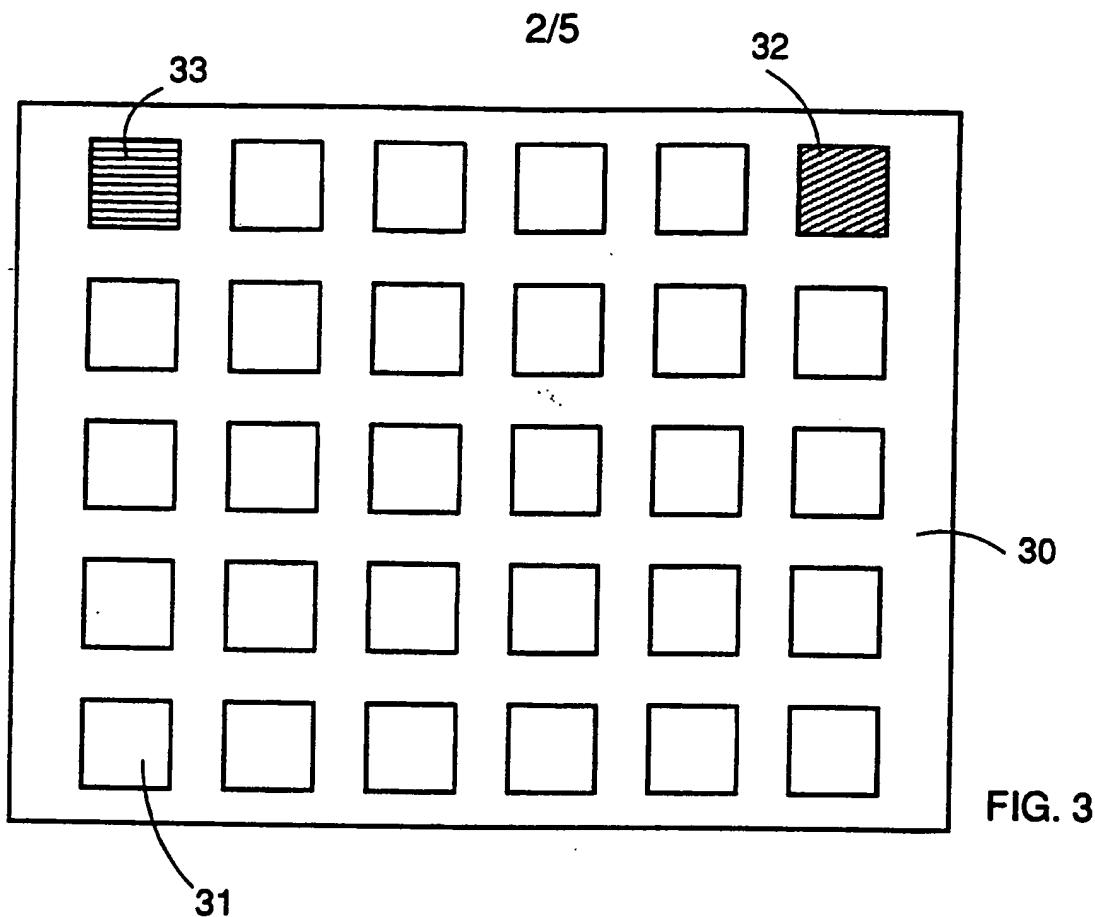


FIG. 2



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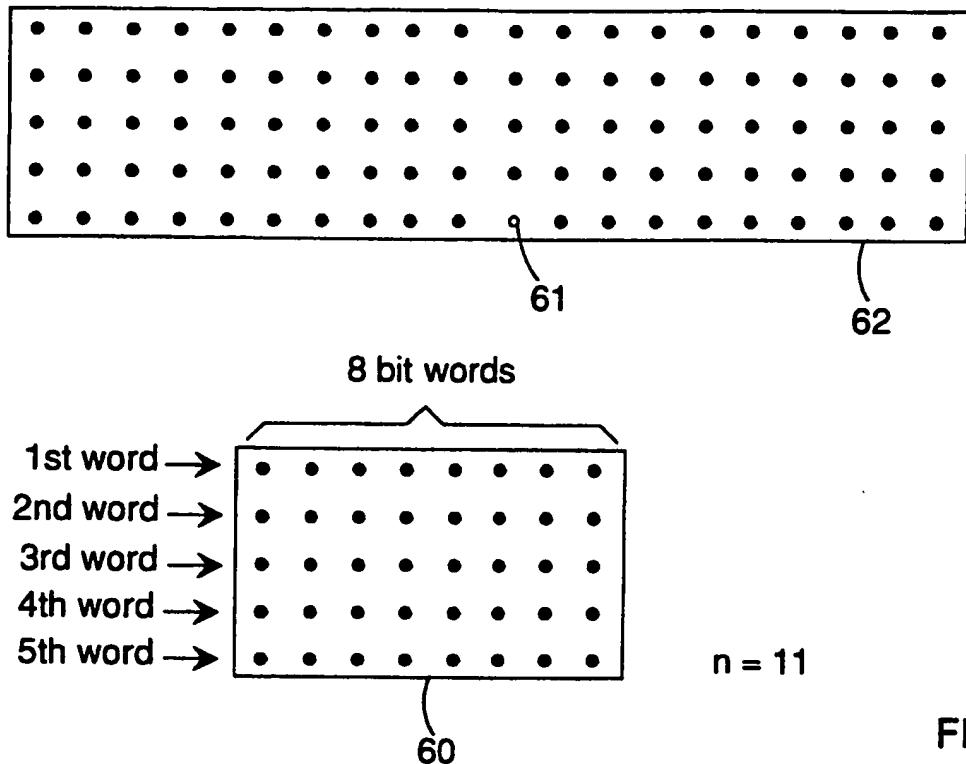


FIG. 6

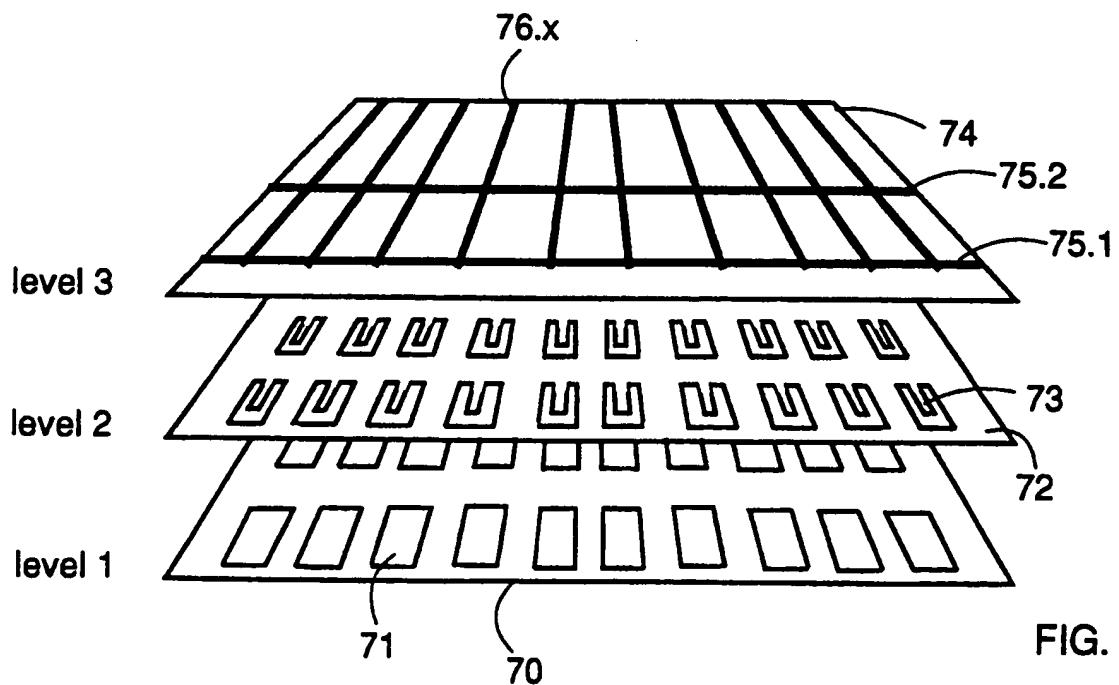
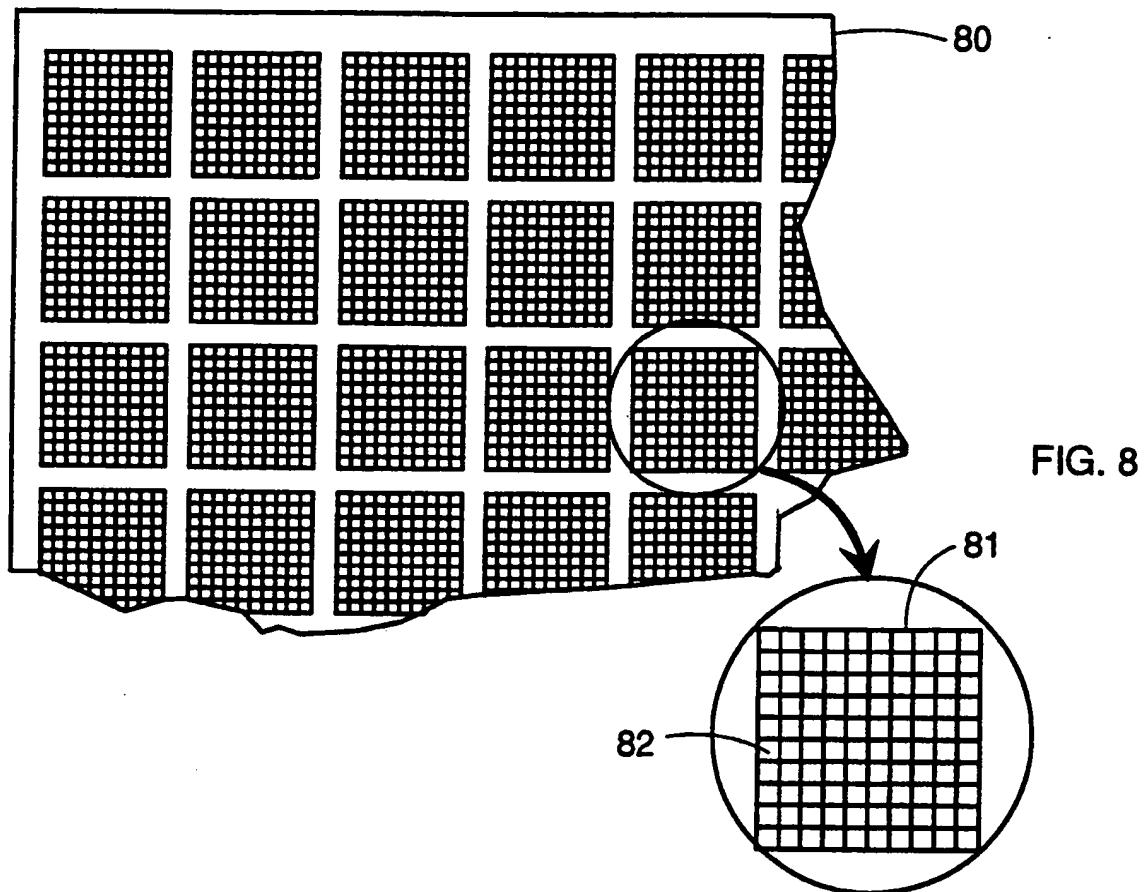
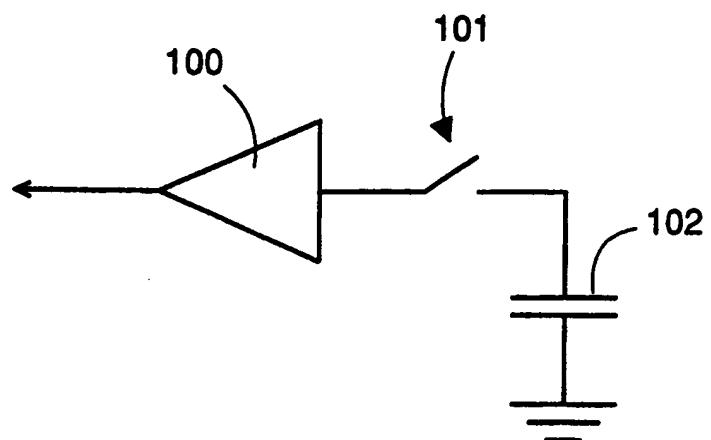
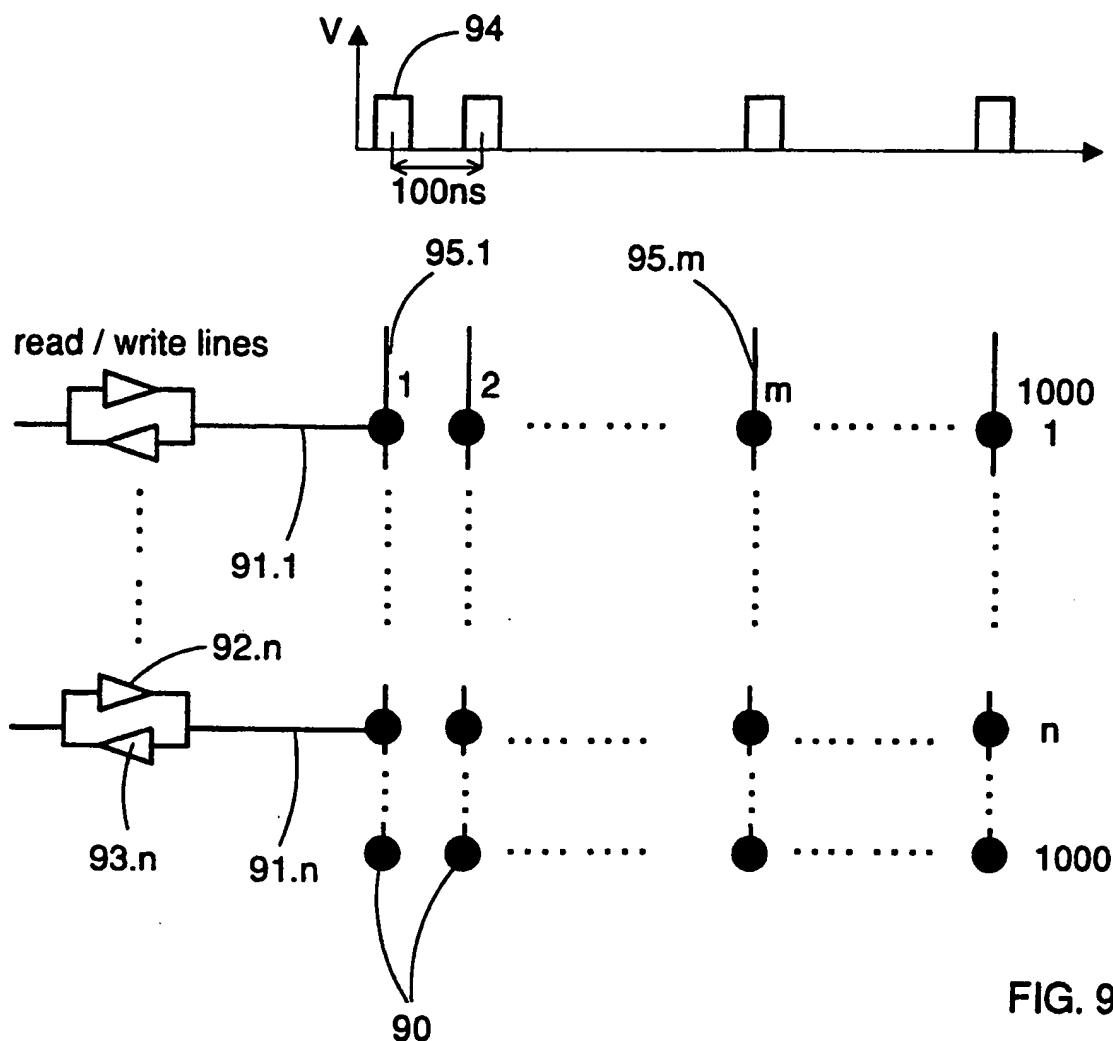


FIG. 7

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INTERNATIONAL SEARCH REPORT

Int. onal Application No.
PCT/IB 95/00594

A. CLASSIFICATION OF SUBJECT MATTER
IPC 6 G11B9/00 G11B19/02 G11B19/04 G11B27/32

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)
IPC 6 G11B G11C

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
Y A	US,A,5 036 490 (KAJIMURA HIROSHI ET AL) 30 July 1991 see column 2, line 66 - column 3, line 11 see column 4, line 50 - line 64 see column 5, line 33 - line 54 see column 6, line 13 - line 19 see column 6, line 46 - column 7, line 10 see column 7, line 25 - line 37 see column 7, line 60 - column 8, line 8; claims 10,11; figures 1,5,6,8-10 ---	1,27 4-6,17, 30
Y A	EP,A,0 600 511 (SONY CORP) 8 June 1994 see page 5, line 9 - page 6, line 37; figure 2 4 5 7A --- -/-	1,27 6,7,10, 28

Further documents are listed in the continuation of box C.

Patent family members are listed in annex.

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- *'T' later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention
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Date of the actual completion of the international search 19 March 1996	Date of mailing of the international search report - 9. 05. 96
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INTERNATIONAL SEARCH REPORT

International Application No
PCT/IB 95/00594

C(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
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